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DATA STRUCTURE

PSEUDO-CODE EXAMPLES

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Note: The Data Structures is a fairly complicated topic with lots of small things to take care of. As such it is quite difficult not to include these small things as they make quite a difference in operating these data structures. I have tried to keep things as simple as possible, leaving difficult things out where possible. However, you can choose for yourself which portions you want discard and which ones to use in your lectures. One example of this could be the "finding of required Node" in the Deletion section of Linked List. The entire "finding" section could be replaced with:

```
Find the required Node
If Node is found
.....
Otherwise
.....
Endif
```

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Stacks

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We need two variables. One for the maximum limit of the Stack (MAX) and the other for keeping track of the current element to be manipulated (Current). We also declare an array for the Stack.

```
MAX = 5
Current = 0
Array Stack[MAX]
```

'This will declare an array with 5 elements.

Inserting a Value onto the Stack (Push)

```
Procedure Push (Value)
If Current is equal to MAX then
    Output an error that the Stack is full
```

'This means no more space for data. Stack is full.

'Display an error message that there is no more space on the Stack.

```
Else
  Add 1 to Current
  Put data in Value into current position
Endif
End Procedure
```

'Go to the next available element in the Stack.
'Put (Push) the data sent in Value parameter in the new element.

Displaying a value from the Stack (Pop)

```
-----
Procedure Pop
if Current is equal to Zero then
  Output error that Stack is empty
Else
  Output value from Stack at Current position
  Subtract 1 from Current
Endif
End Procedure
```

'If there are no values in the Stack.
'Nothing to output.
'Output (POP) a value in the Current element (LIFO).

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QUEUES

Prerequisites:

```
End = 0
MAX = 5

Array Queue[MAX]
```

To track the End element of the Queue.
Maximum number of elements (or whatever value you like as the number of elements in the Queue)
Declares an array of 5 elements for the Queue.

Adding a Value

```
-----
Procedure Insert(Value)
If End is equal to MAX then
  Output an error message that the Queue is full
Otherwise
  Add 1 to End
  Store data in Value into Queue at End position
Endif
End Procedure
```

'No more space in Queue
'Go to next available element in Queue

Displaying value from Queue

Procedure Display

If End is equal to Zero then

'No values to display

 Output an error message that the Queue is empty

Otherwise

'Always display first value (FIFO)

 Display value at position Queue[1]

'Looping through all elements with values

 Start Loop from position 1 to End position

'Moving data is step on its left to fill the gap in the 1 position

 Copy data from Next element to Current element

 Move to Next element

 End Loop

EndIf

End Procedure

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Linked List

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Declaring a Linked List

Integer : Data

Pointer: Next

Declaring variables

Pointer : Start

Start = NULL

'Nothing in the Linked List

Creating a New Node with above structure.

Procedure Create(Value)

If there is enough memory available then

'Only create a Node if memory is available

 Create New Node in memory

'Creating a New Node for the Linked List

 Point its Next pointer to NULL

'Making a termination point for the Node

 If this is the first Node in the List then

'There must be a starting point for the Linked List

 Point Start to the Node

'Point Start to New Node only if it is the very first one

<pre> Otherwise Go to memory location pointed by Start pointer Start Loop and run until Next points to NULL Go to the Node pointed by Next pointer End Loop Point Next of last Node to the newly created Node Put data in Value in new Node's Data Endif Otherwise Output error message that there is no more memory available Endif End Procedure </pre>	<pre> 'If other Nodes exist then find the last Node in the List 'Begin at the Start (first Node) 'Check each Node until Next pointer points to NULL 'Move from Node to Node 'When last Node found, point its Next pointer to New Node 'Put the required data (Value) in New Node's Data variable 'Out of memory. Can't create a New Node </pre>
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Deleting a Node from the Linked List (assuming values in Data are unique)

Procedure Delete(Value)

Pointer : Found

Found = NULL

If there are no Nodes in the List then

Output error message that can't delete a Node

Otherwise

Go to memory location pointed by Start pointer

Start Loop and run until Next points to NULL

If value in Data of current Node is equal to Value then

Found = current Node

Exit Loop

Endif

Go to Node pointed by Next pointer of current Node

End Loop

If the required Node is the first Node in the List then

Point the Start pointer to NULL

Otherwise

Point Next of previous Node to the Node pointed by Next of Found 'Node deleted (Previous.Next = Found.Next)

Endif

Endif

End Procedure

'Pointer for finding required Node to Delete

'Assuming required Node not found

'Begin at the Start (first Node)

'Check each Node until Next pointer points to NULL

'If we find our required Node

'Mark the current Node with Found pointer

'No need to finish the loop

'Node deleted

Displaying the entire Linked List

Procedure Display

If Start points to NULL then

'List is empty

Output error message that the List is empty

Otherwise

Start Loop at Start position and run until Next points to NULL

Output current Node's Data

Go to Node pointed by Next pointer of current Node

End Loop

Endif

End Procedure

Binary Trees

Structure of Binary Tree Node

Integer : Data

Pointer : Left

Pointer : Right

Pointer : Start

Pointer : Current

Start = NULL

Current = NULL

Insertion in a Binary Tree

Procedure Insertion(Value)

Create a New Node

Point its Left and Right pointers to NULL

If Start is pointing to NULL then

Point Start to the New Node

'Tree was empty so insert at Start

Otherwise

Current = Start

'Search for appropriate position

Start Loop at Current and run until Left and Right pointers are NULL

```
If data in Value is less than or equal to Data in Current Node then
  Go to the left of the Current Node
Otherwise
  Go to the right of the Current Node
Endif
If data in New Node is less than data in Current Node then
  Point the Left pointer of Current Node to the New Node
Otherwise
  Point the Right pointer of Current Node to the New Node
Endif
Insert data in Value in the New Node
End Loop
Endif
End Procedure
```

Traversals of Binary Tree

Procedure InOrder(Current)

```
If Left pointer of Current is not NULL then
  Call InOrder with Left pointer of Current
Endif
```

'This is a recursive call from InOrder to itself

Output Data in Current Node

```
If Right pointer of Current is not NULL then
  Call InOrder with Right pointer of Current
Endif
```

'This is a recursive call from InOrder to itself

End Procedure

Procedure PreOrder(Current)

Output Data in Current Node

```
If Left pointer of Current is not NULL then
  Call InOrder with Left pointer of Current
Endif
```

'This is a recursive call from PreOrder to itself

```
If Right pointer of Current is not NULL then
  Call InOrder with Right pointer of Current
Endif
```

'This is a recursive call from PreOrder to itself

End Procedure

Procedure PostOrder(Current)

If Left pointer of Current is not NULL then

 Call InOrder with Left pointer of Current

Endif

If Right pointer of Current is not NULL then

 Call InOrder with Right pointer of Current

Endif

Output Data in Current Node

End Procedure

'This is a recursive call from PostOrder to itself

'This is a recursive call from PostOrder to itself

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